Wireless and Communication in the Internet of Things
BLE Connections

Pat Pannuto, UC San Diego

ppannuto@ucsd.edu

Today's Goals

- Explore how connections work
  - What does the link layer look like?
  - How do higher layers interact to share data?

- Investigate network questions about connections

- Overview of additions in BLE 5.0

Useful documentation

- [5.2 specification] [4.2 specification] (link to PDF download)

- https://www.novelbits.io/deep-dive-ble-packets-events/

- https://www.bluetooth.com/bluetooth-resources/understanding-reliability-in-bluetooth-technology/

- Thinking about BLE connection data transfer rates
  1. https://punchthrough.com/maximizing-ble-throughput-on-ios-and-android/
  2. https://punchthrough.com/maximizing-ble-throughput-part-2-use-larger-att-mtu-2/
  3. https://punchthrough.com/maximizing-ble-throughput-part-3-data-length-extension-dle-2/

Outline

- **Connection PHY and Link Layer**

- Connections as Networks

- GATT


- BLE 5

Overview of connections

- Connections are for bi-directional communication with higher throughput than advertisements

- Simple view
  - A peripheral is either advertising or in a connection
  - A central is scanning and in one or more connections
  - (Remember: actually false, devices can have many roles simultaneously)

- While in a connection <u>both</u> devices act like servers
  - Either device can read/write fields available on the other device

# Once a connection is established, BLE has more options

- LE 1M [default]
  - 1 Msys/s   [n.b. BLE encodes 1 bit / symbol, so this is also 1 Mbit/s]
- LE 2M
  - 2 Msys/s
- LE Coded
  - 1 Msys/s + FEC
    - S = 2, ~2x range, ½ effective data rate – 500 Kbit/s **goodput**
    - S = 8, ~4x range, 125 Kbits/s **goodput**

Frequency Hopping Spread Spectrum (FHSS)

- Recall: Each BLE channel is 2 MHz wide, 40 channels,
  - 3 are used for advertising, remaining 37 are for connections
  - Frequency hopping: $f_{n+1} = (f_n + \mathtt{hop})\ \mathtt{mod}\ 37$

- Which exact channels are used and in what order might vary
  - "Adaptive Frequency Hopping" avoids bad channels
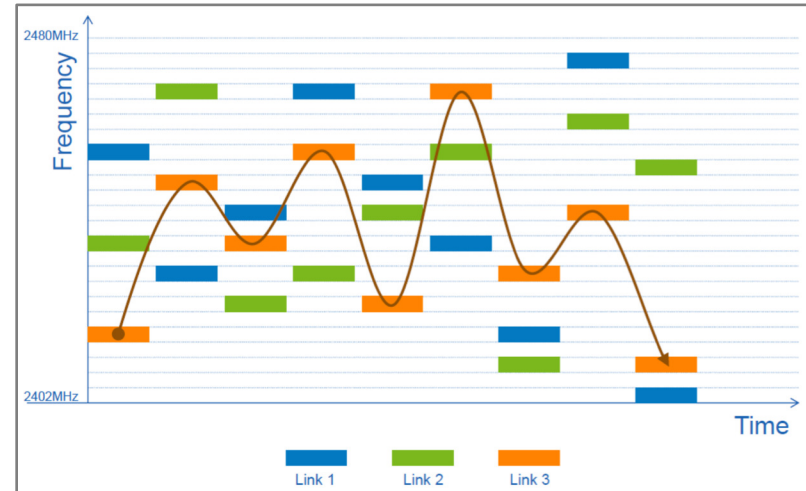
Connection timeline (in picture form)

CC BY-NC-ND Pat Pannuto – Content developed in collaboration with Branden Ghena
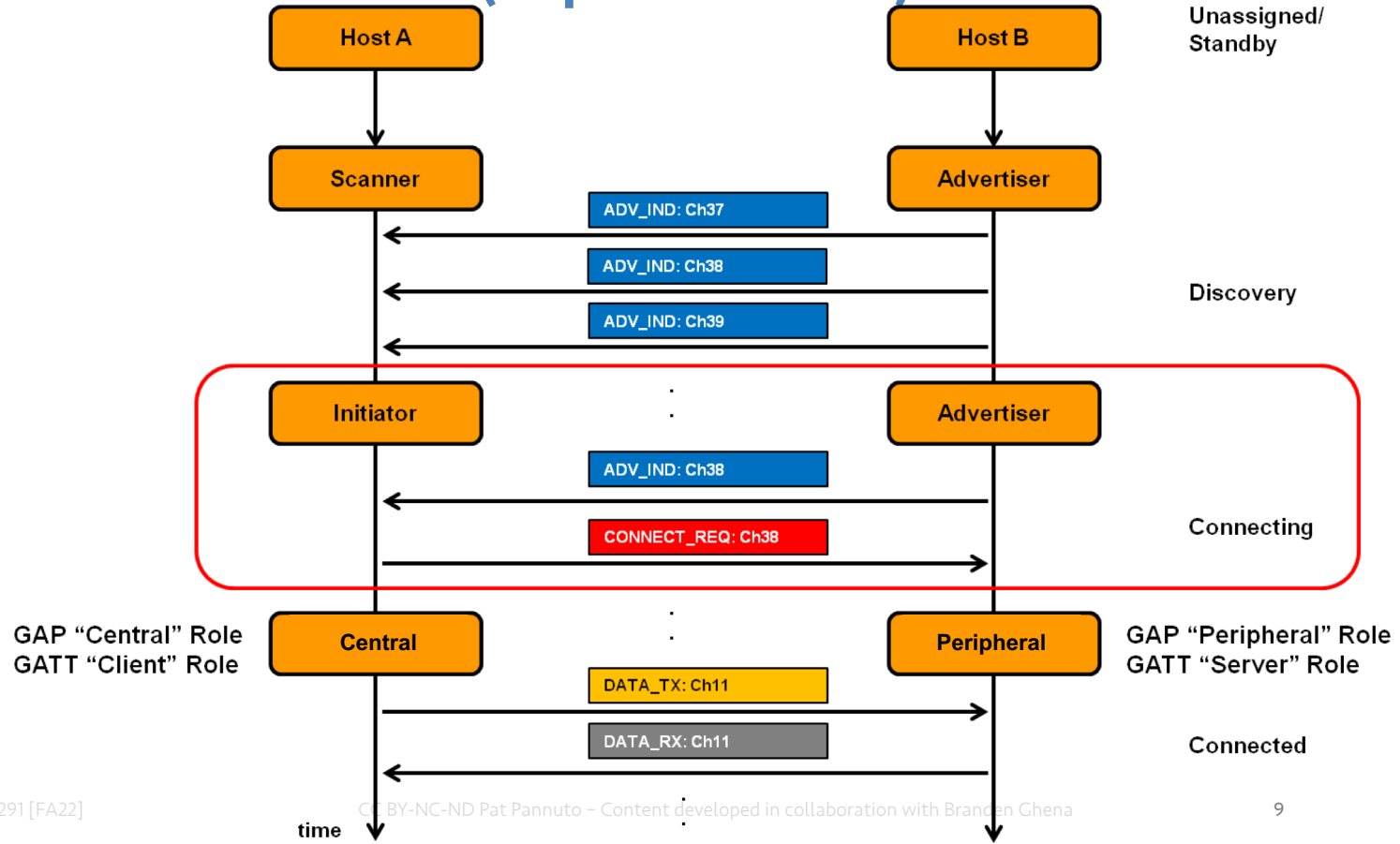
Connection timeline

- Initiating a connection
  - Peripheral is sending broadcast advertisements
  - Central is scanning and receives an advertisement
  - Central sends connection request

- During a connection
  - Central sends a packet each "connection interval"
  - Peripheral immediately responds with a packet
  - Multiple packets may be exchanged this way until done
  - Repeat at next connection interval

- Ending a connection
  - Either device sends termination packet
  - Timeout occurs on either device

Connection request packet

- Scanner waits until it sees an advertisement from a device it wants to connect to
  - Requesting a connection (in higher layers) just starts this search process

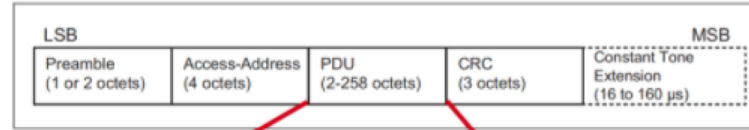- Sends connection request payload in response to advertisement

Figure 2.1: Link Layer packet format for the LE Uncoded PHYs
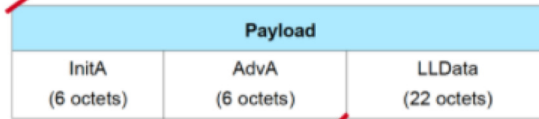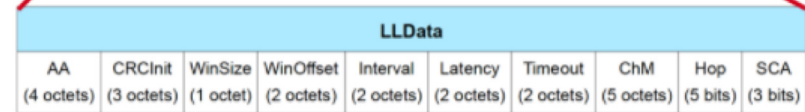
Figure 2.4: Advertising physical channel PDU

Figure 2.12: CONNECT_IND and AUX_CONNECT_REQ PDU Payload

Request parameters

| LLData									
AA (4 octets)	CRCInit (3 octets)	WinSize (1 octet)	WinOffset (2 octets)	Interval (2 octets)	Latency (2 octets)	Timeout (2 octets)	ChM (5 octets)	Hop (5 bits)	SCA (3 bits)

- Specifies parameters of the connection
  - Peripheral must either agree or totally reject the connection
  - Peripheral can later propose a change to the connection parameters

- Interval is how frequently connection events occur
- Timeout is how long since hearing from a device before the connection breaks

- Channel Map and Hop have to do with FHSS pattern

The Central schedules the first connection event

- WinSize and WinOffset specify start of the first connection event
  - Places an "anchor" point that defines the TDMA schedule for this device
  - Interval specifies duration between connection events starting at "anchor"
  - Allows Central to place this connection, avoiding its other connections
    - Before first response from peripheral, timeouts are faster
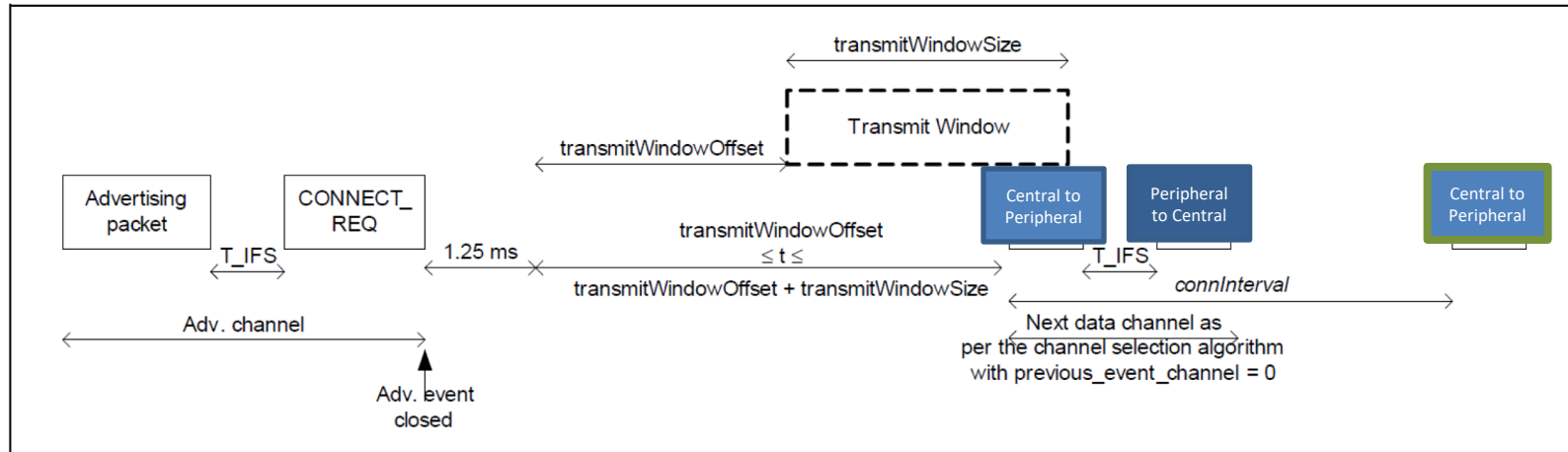
Steady-state connection timing
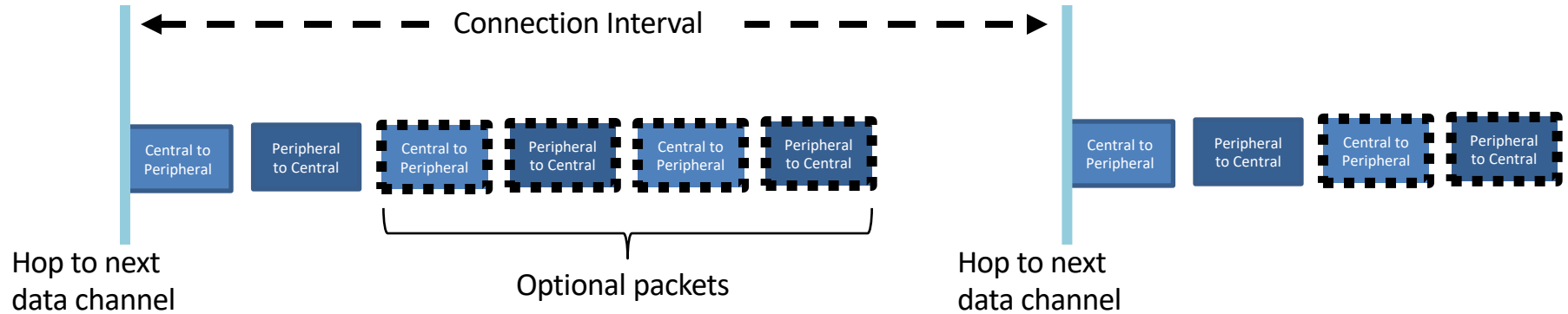
- **Some data can be exchanged at each interval**
  - Might just be acknowledgements
  - Additional packets can be sent if there is a lot to transmit
  - Each interval is on the next channel in the hopping sequence

- **Peripheral can skip a number of intervals to save energy**
  - Defined by the Latency connection request parameter

Connection packet layering

LSB			MSB
Preamble (1 octet)	Access Address (4 octets)	PDU (2 to 257 octets)	CRC (3 octets)

Figure 2.1: Link Layer packet format

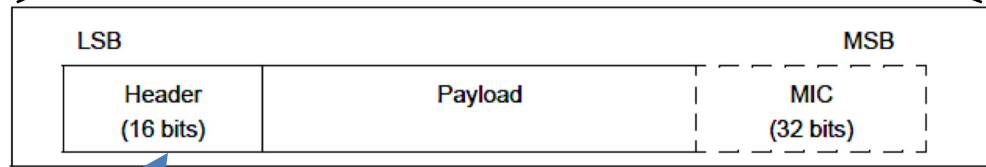| LSB		MSB
Header (16 bits)	Payload	MIC (32 bits)

Figure 2.12: Data Channel PDU

Once we have an established connection…

How do we manage this link?

Link Layer Header

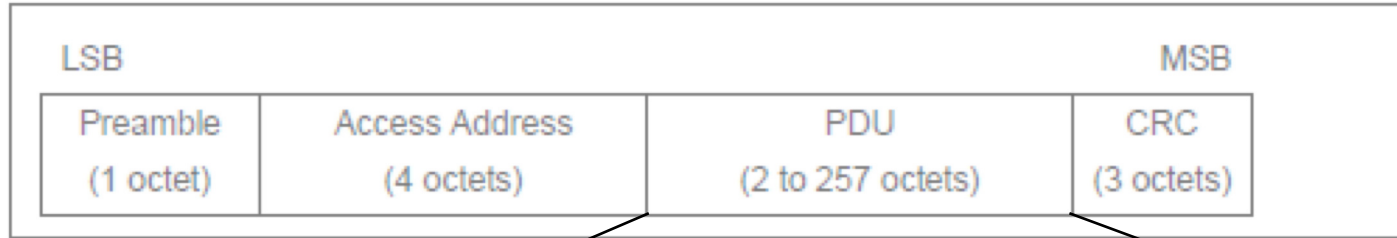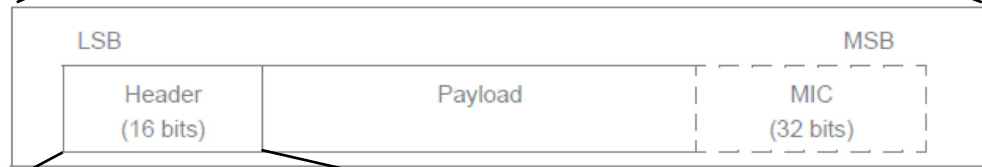Figure 2.1: Link Layer packet format

[Next Expected] Sequence number – ACK/dupe protocol

Figure 2.12: Data Channel PDU
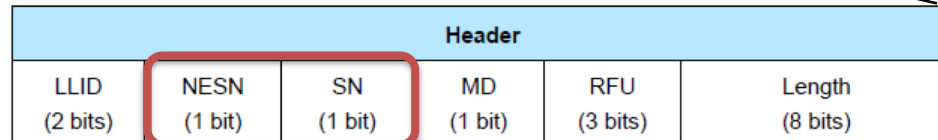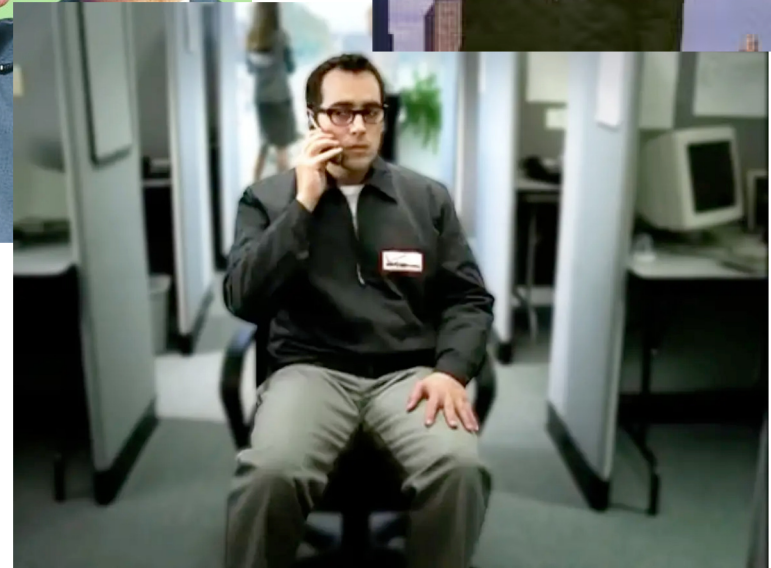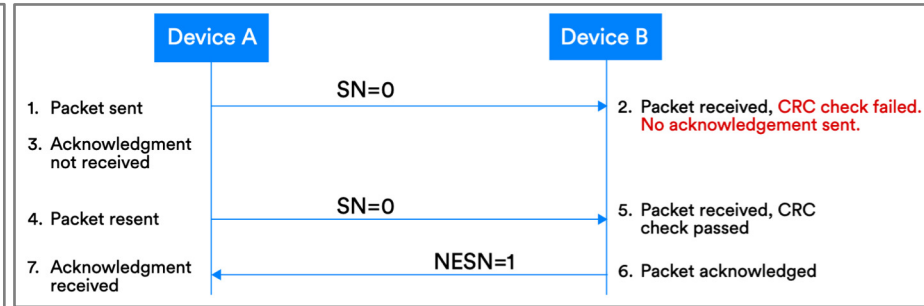
Figure 2.13: Data channel PDU header

Acknowledgements ensure your message was received

BLE Acknowledgement (ACK) protocol
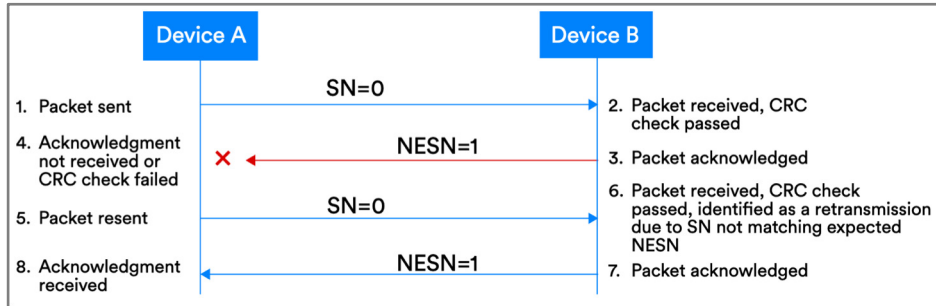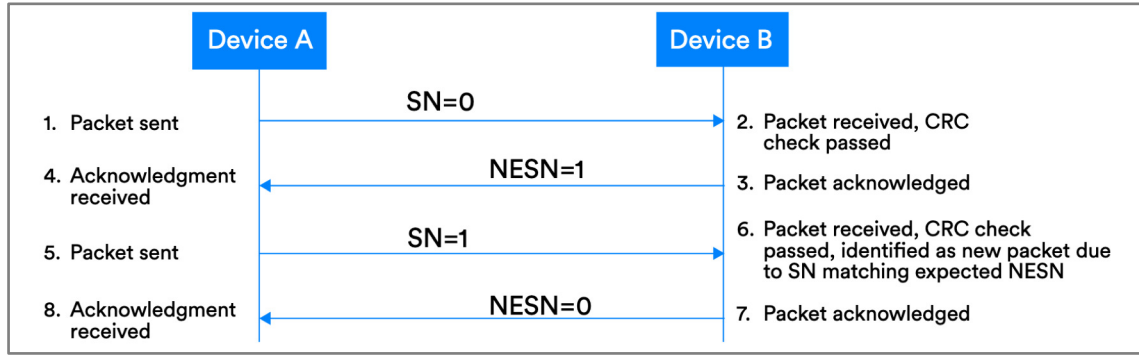
See §4.2.3 of Reliability Doc for more

Question: Are ACKs sent for advertisements or connections (or both)?

Link Layer Header
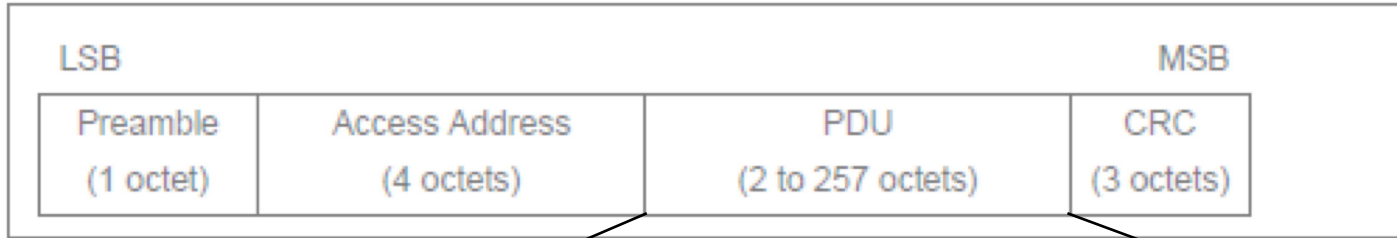
Figure 2.1: Link Layer packet format

LLID – link layer ID
* Empty data / Fragment
* Data
* Control

MD – more data (continues connection event)

RFU – Reserved for Future Use
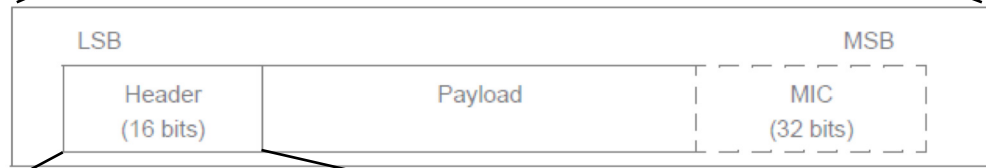
Figure 2.12: Data Channel PDU
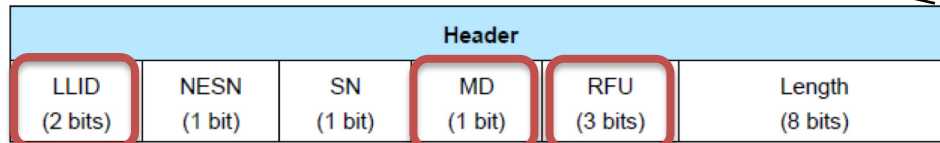
Figure 2.13: Data channel PDU header

Link Layer Header
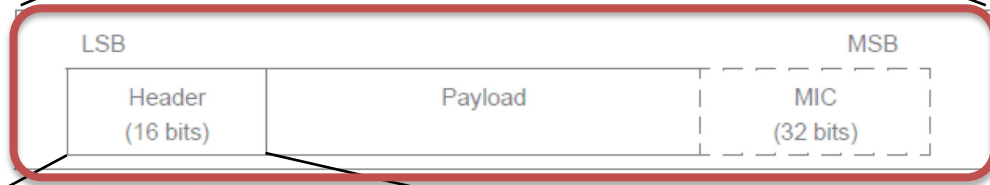
Figure 2.1: Link Layer packet format

Question:
What is the maximum possible payload size?

Figure 2.12: Data Channel PDU

Figure 2.13: Data channel PDU header

Parsing packet headers

LSB			MSB
Preamble (1 octet)	Access Address (4 octets)	PDU (2 to 257 octets)	CRC (3 octets)

Figure 2.1: Link Layer packet format

LSB		MSB
Header (16 bits)	Payload	MIC (32 bits)
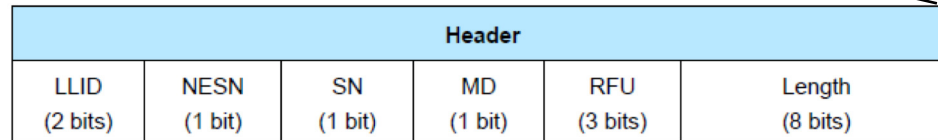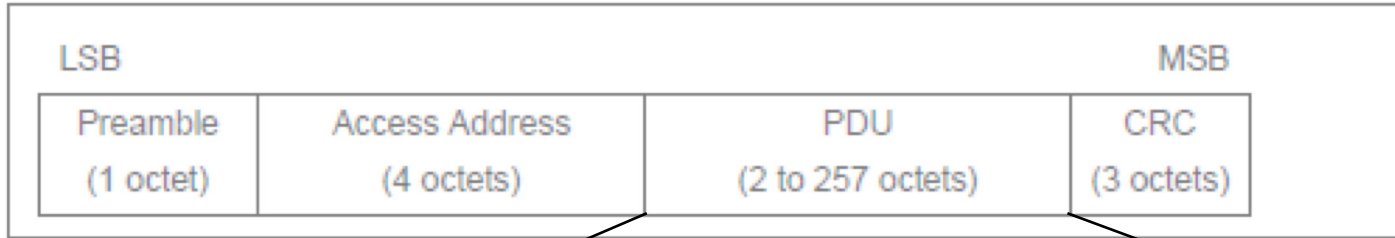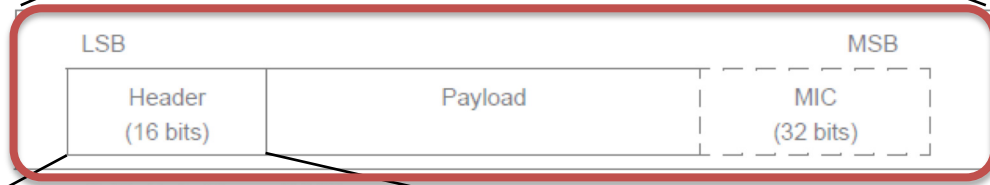
Figure 2.12: Data Channel PDU

Question: Could you have answered that without the header information?

MIC = Message Identification Code
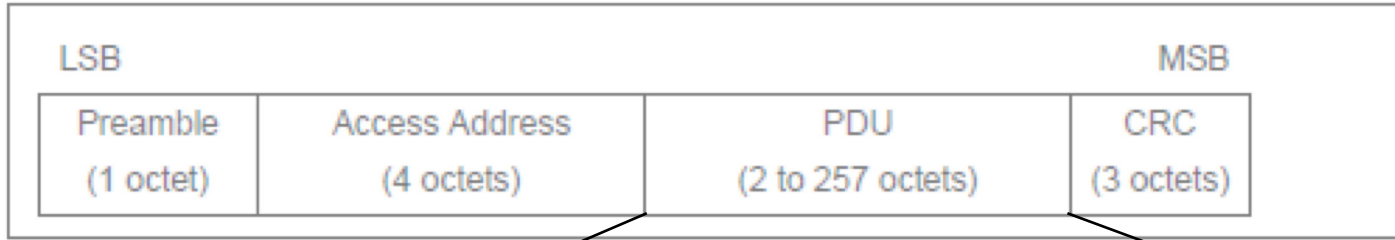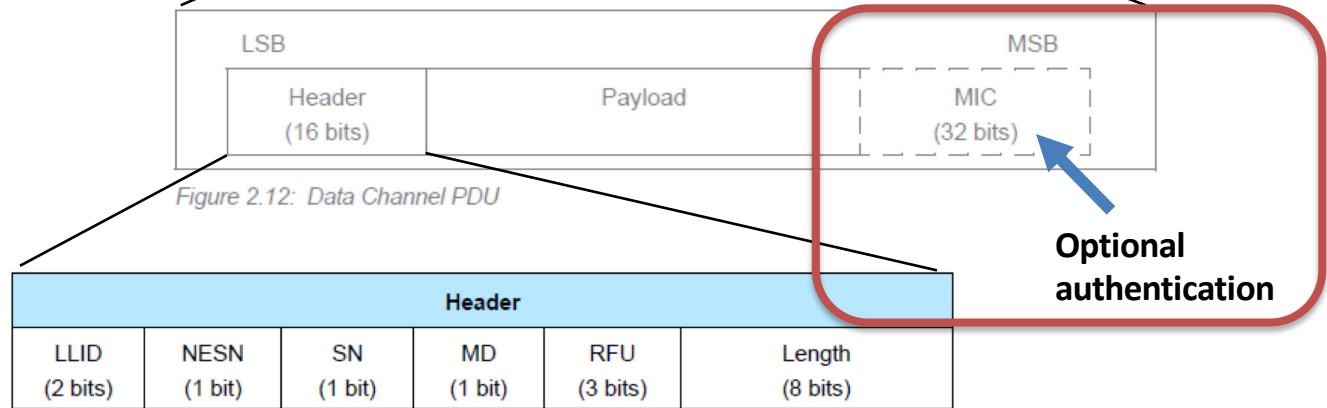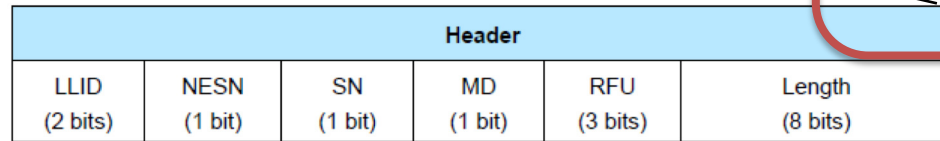
LSB			MSB
Preamble (1 octet)	Access Address (4 octets)	PDU (2 to 257 octets)	CRC (3 octets)

Figure 2.1: Link Layer packet format

LSB		MSB
Header (16 bits)	Payload	MIC (32 bits)
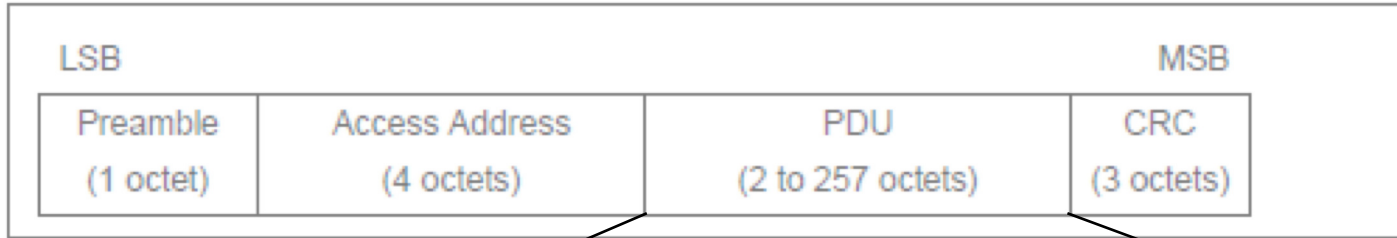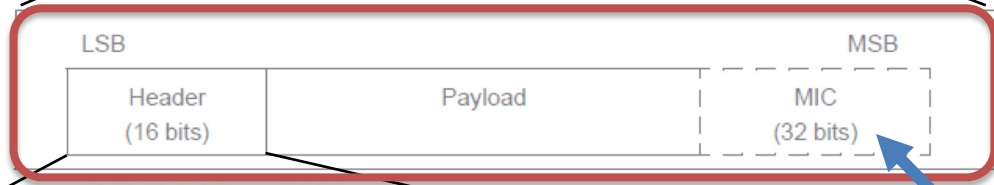
Figure 2.12: Data Channel PDU

Optional authentication

Header					
LLID (2 bits)	NESN (1 bit)	SN (1 bit)	MD (1 bit)	RFU (3 bits)	Length (8 bits)

Figure 2.13: Data channel PDU header

Parsing Packet Headers

Preamble (1 octet)	Access Address (4 octets)	PDU (2 to 257 octets)	CRC (3 octets)

LSB   MSB

Figure 2.1: Link Layer packet format

LSB   MSB

Header (16 bits)	Payload	MIC (32 bits)

Figure 2.12: Data Channel PDU

Question:
Is the MIC part of the payload, or after it?
How do you know?

Optional authentication

Header					
LLID (2 bits)	NESN (1 bit)	SN (1 bit)	MD (1 bit)	RFU (3 bits)	Length (8 bits)

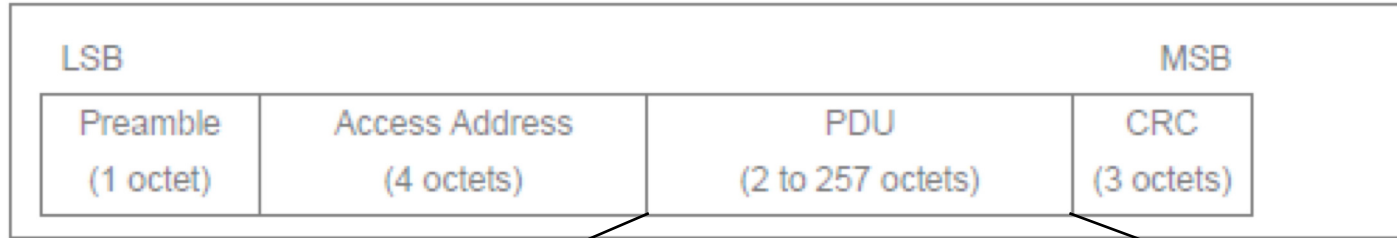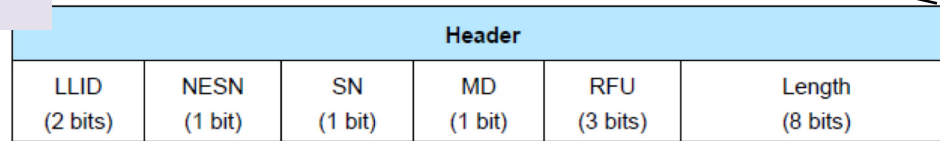Figure 2.13: Data channel PDU header

Sometimes, you still pay for what you don't use
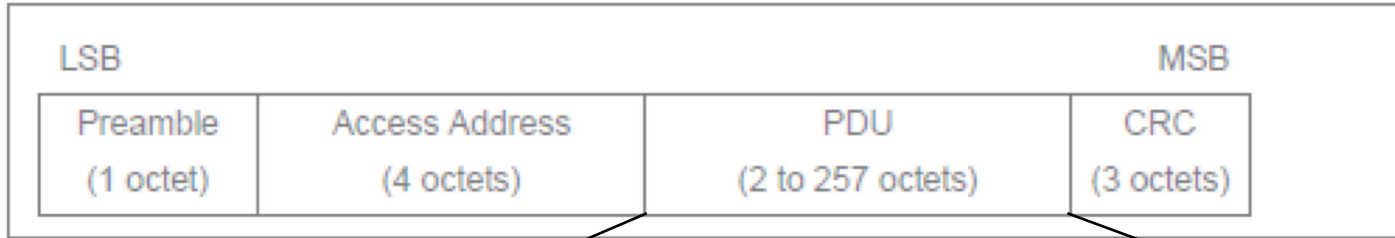
Figure 2.1: Link Layer packet format
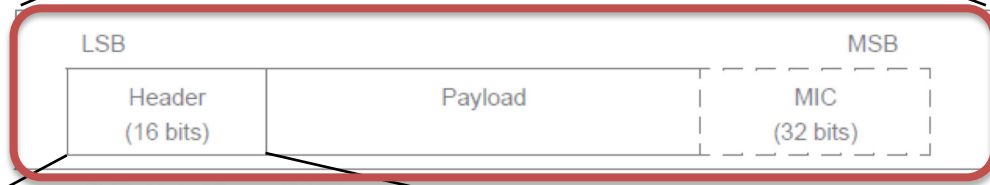
Question:
What is the maximum possible <u>payload</u> size?

Figure 2.12: Data Channel PDU
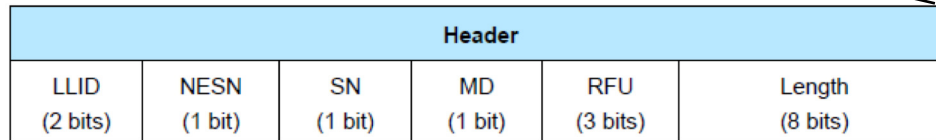
Figure 2.13: Data channel PDU header

"Control" and "Data" separation — very common theme in networking

- LLID – link layer ID
  - Empty data / Fragment
  - Data
  - Control
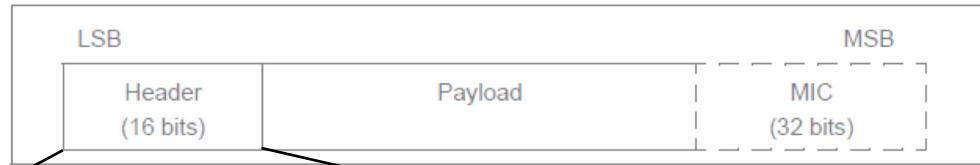
Figure 2.12: Data Channel PDU
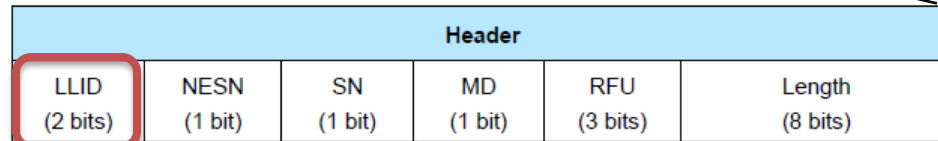
Figure 2.13: Data channel PDU header

Control payloads

Examples:

- Request update to connection parameters like interval (by peripheral)

- Begin encrypting communication

- Terminate a connection

Notice: Control payloads can all be handled by the link layer – not passed up!

Opcode	Control PDU Name
0x00	LL_CONNECTION_UPDATE_REQ
0x01	LL_CHANNEL_MAP_REQ
0x02	LL_TERMINATE_IND
0x03	LL_ENC_REQ
0x04	LL_ENC_RSP
0x05	LL_START_ENC_REQ
0x06	LL_START_ENC_RSP
0x07	LL_UNKNOWN_RSP
0x08	LL_FEATURE_REQ
0x09	LL_FEATURE_RSP
0x0A	LL_PAUSE_ENC_REQ
0x0B	LL_PAUSE_ENC_RSP
0x0C	LL_VERSION_IND
0x0D	LL_REJECT_IND
0x0E	LL_SLAVE_FEATURE_REQ
0x0F	LL_CONNECTION_PARAM_REQ
0x10	LL_CONNECTION_PARAM_RSP
0x11	LL_REJECT_IND_EXT
0x12	LL_PING_REQ
0x13	LL_PING_RSP
0x14	LL_LENGTH_REQ
0x15	LL_LENGTH_RSP

What's an example of link layer control?
(Oh, and by the way, I've been lying little bit about PDUs)

- In BLE, the maximum PDU is 27 bytes
  - i.e., connection and advertisements are the same length

- In BLE 4.2, we get the Data Length Extension (DLE) feature
  - Kinda does what it sounds like, allows for PDUs up to 257 bytes (~10x!)

Central — Peripheral

LL_FEATURE_REQ
LE Data Packet Length Extension = Enabled

LL_FEATURE_RSP
LE Data Packet Length Extension = Enabled

LL_LENGTH_REQ
connMaxRxOctets = 251
connMaxRxTime = 2120 μs
connMaxTxOctets = 251
connMaxTxTime = 2120 μs

LL_LENGTH_RSP
connMaxRxOctets = 251
connMaxRxTime = 2120 μs
connMaxTxOctets = 127
connMaxTxTime = 1128 μs

Master determines that:
MaxRxOctets = 127
MaxRxTime = 1128 μs
MaxTxOctets = 251
MaxTxTime = 2120 μs

Slave determines that:
MaxRxOctets = 251
MaxRxTime = 2120 μs
MaxTxOctets = 127
MaxTxTime = 1128 μs

PunchThrough

https://punchthrough.com/maximizing-ble-throughput-part-3-data-length-extension-dle-2/

LSB			MSB
Preamble (1 octet)	Access Address (4 octets)	PDU (2 to 257 octets)	CRC (3 octets)

Figure 2.1: Link Layer packet format

"Control" and "Data" separation — very common theme in networking

- LLID – link layer ID
  - Empty data / Fragment
  - Data
  - ~~Control~~

Figure 2.12: Data Channel PDU
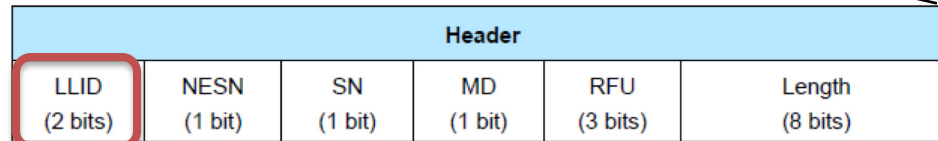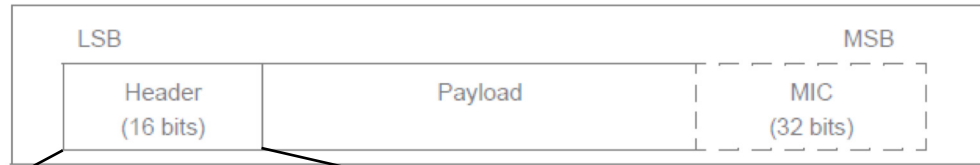
Figure 2.13: Data channel PDU header
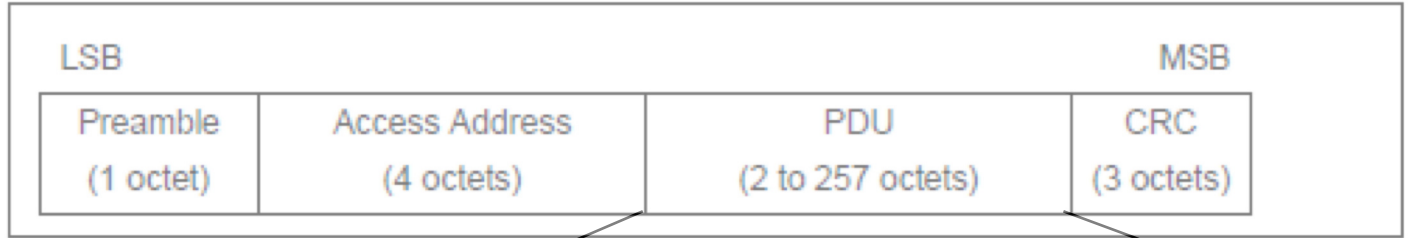
Data payloads

Figure 2.1: Link Layer packet format

| LSB			MSB
Preamble (1 octet)	Access Address (4 octets)	PDU (2 to 257 octets)	CRC (3 octets)

Figure 2.12: Data Channel PDU

LSB		MSB
Header (16 bits)	Payload	MIC (32 bits)

Length (16 bits)	Channel ID (16 bits)	L2CAP SDU Length* (16 bits)	Information Payload (0 to **???** octets)

Data payloads

LSB ... MSB

| Preamble (1 octet) | Access Address (4 octets) | PDU (2 to 257 octets) | CRC (3 octets) |
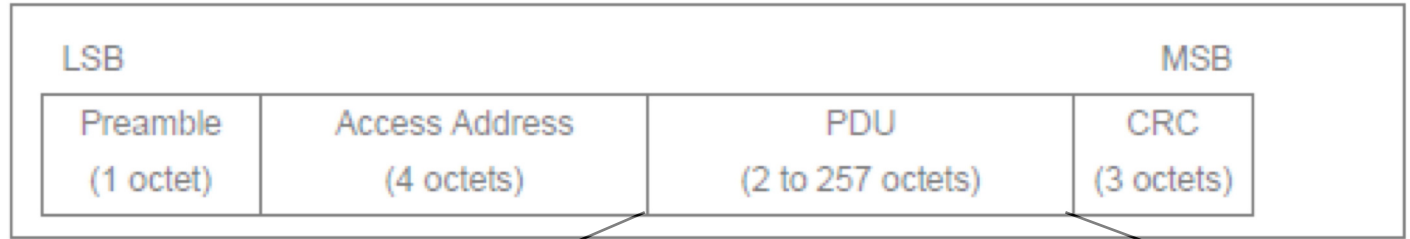
Figure 2.1: Link Layer packet format

SDU – Service Data Unit
- Logical packet that may be split into several physical packets
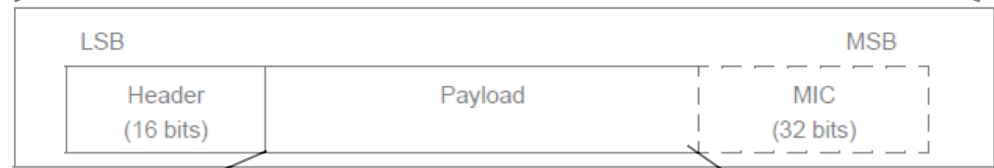
Channel ID
- Destination for data
- ATTributes or Security
- Can also make custom endpoints

LSB ... MSB

| Header (16 bits) | Payload | MIC (32 bits) |

Figure 2.12: Data Channel PDU

| Length (16 bits) | Channel ID (16 bits) | L2CAP SDU Length* (16 bits) | Information Payload (0 to 249 octets) |
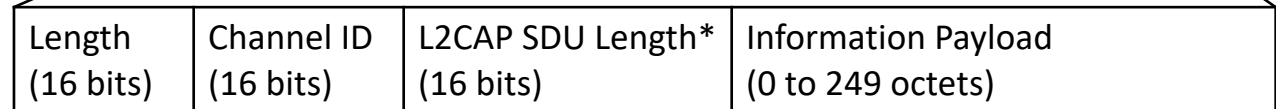
Zooming back out, summarizing what's sent in a BLE connection

- Control packets
  - Version
    - Bluetooth version, Company ID of Software Stack
  - Features
    - Various connection features supported by the device
    - Encryption, Ping, Channel selection algorithms, Various 5.0 features
  - Etc.
    - Extended packet length, Various 5.0 features (change PHY)
    - Various "procedures" for setting up changes in the connection

- Data packets
  - Attribute discovery
  - Attribute reading

Question: Which kind of packets do you think are sent more often, data or control?

- How might device type and activity play a role in this?

Very briefly: Security in connections

Either side can request secure connection [central -> Pairing; peripheral -> Security Req]

Pairing Request Opcode (0x01)	IO Capability	OOB Data Flag	AuthReq	Max Encryption Key Size	Initiator Key Distribution	Responder Key Distribution

Pairing Request Command Format

Pairing Response Opcode (0x02)	IO Capability	OOB Data Flag	AuthReq	Max Encryption Key Size	Initiator Key Distribution	Responder Key Distribution

Pairing Response Command Format

Security Request Opcode (0x0B)	AuthReq

Peripheral Security Request Command Format

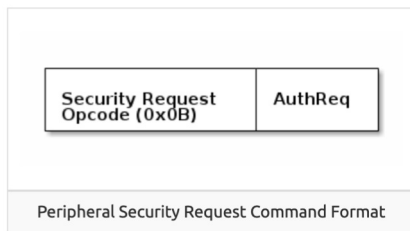| Association Models	Conditions for Use
Just Works	This model is used when there is no Out of Band (OOB) data available and neither device has input capabilities.
Passkey Entry	This model is used when there is no OOB data available, at least one device has the I/O capability to enter a passkey, and the other device has the capability to display a passkey.
Numeric Comparison	This model is used if both devices support Secure Connections, can display a yes or no message and have some input capability.
Out of Band	If OOB data is available on either device, this model will be chosen.

"Bonding" saves pairing
info between connections

- Excellent TI resource here for more details:
  https://dev.ti.com/tirex/explore/node?node=AOPOY.GDApakIOYjiwoY6A__pTTHBmu__LATEST

Very Briefly: Security in Addresses

- All devices have a unique, hardware "Identity Address", but may not share it
  - Public Address - Use the Identity Address. (The BDA never changes.)
  - Random Static Address - Generate a random address per power cycle. The address cannot be regenerated at any other time. Can be used as an Identity Address.
  - Resolvable Private Address (RPA) - Generate a random address with a given time interval. Generated using Identity Resolving Key (IRK) which can also be used by trusted peers (bonded devices) to resolve the Random Address to the Identity Address.
  - Non-resolvable Private Address - Generate a random address with a given time interval. Generated randomly. Cannot be resolved to an Identity Address.
- N.b., Nishant is a local expert in BLE address monitoring

Ending connections

- Termination control packet

- Timeout parameter from connection parameters
  - Human-based devices may just wander away from each other
  - Or be shut off, or reprogrammed, etc.

Break + Check your understanding

- What makes connections more reliable than advertisements?

- How does a Peripheral send data to a Central in a connection?

- Why is a "termination" control packet useful?

Break + Check your understanding

- What makes connections more reliable than advertisements?
  - Each packet sent is acknowledged or resent
  - TDMA schedule and FHSS means less likely to collide

- How does a Peripheral send data to a Central in a connection?
  - Wait until next connection interval
  - Respond to packet from Central with a data payload

- Why is a "termination" control packet useful?
  - Timeout could delay for a while
  - Enables re-connection if desired

Outline

- Connection Link Layer

- **Connections as Networks**

- GATT

- BLE 5

Questions about how a connection "network" works

- How does TDMA MAC for connections work?
  - Implies a schedule
  - Implies synchronization

- How many devices can be in a network?

- How much throughput can a device have?

How is the TDMA schedule created/managed?

- Only the central needs to know the whole schedule
  - It controls interactions with the peripheral(s) it is connected to

- Anchor point and connection interval determines the schedule for a specific device
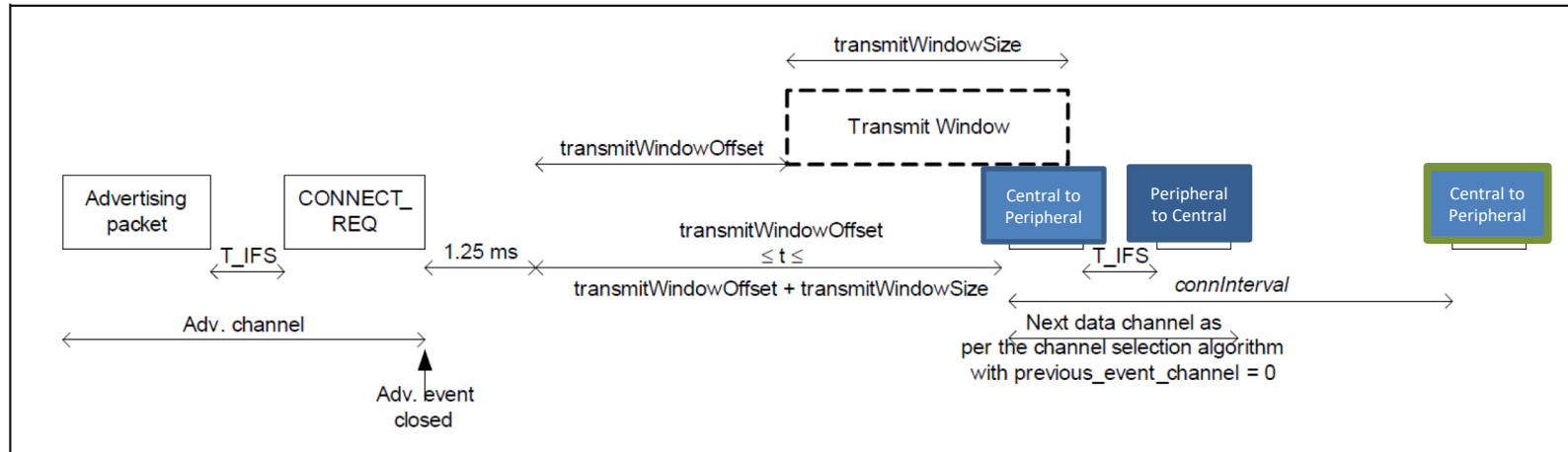
How is synchronization managed?

- Central sends first packet at each connection interval
  - So peripheral must be synchronized to central
  - Resynchronization can occur on each received packet

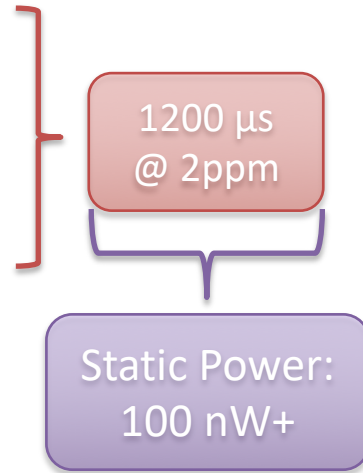- Specification describes how a peripheral must widen its listening window based on "Source Clock Accuracy"

$$windowWidening = \left(\frac{centralSCA + peripheralSCA}{1000000}\right) * timeSinceLastAnchor$$

Clock drift is an energy burden due to large guard bands and energy cost of precise timekeeping

Device time

Guard

300 s

Wall time

1200 µs
@ 2ppm

Static Power:
100 nW+
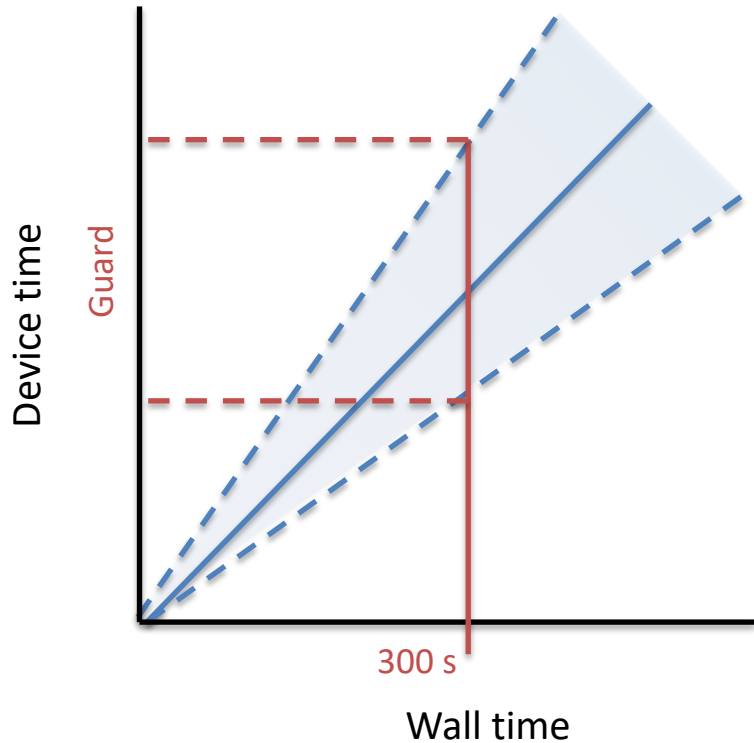
CC2520 Radio + AM18xx RTC:
300s * 2ppm * 2 = 1.2 ms guard
1.2 ms * 60 mW = 36 µJ (RF loss)
300s * 150nW = 45 µJ (RTC loss)
≈ 100 µJ loss / packet @ 300 s

(CC2520 is a 15.4 radio, but the same idea applies)

Max connection interval in BLE is 4 ms

How many devices can be connected?

- We can schedule with granularity of *at least* 1.25 ms (offset steps)
  - That's at least 800 devices per second

- Intervals go up to 4 seconds, so multiply by four
  - That's 3200 devices per max interval

- Plus central can skip intervals on occasion without dropping the connection
  - And really the offset defines a window. More granularity is available

- Answer: thousands of devices
  - Although each is sending minimum-sized packets each interval

How many devices can be connected in the real world?

- The limit is much much lower on real devices
  - Example: Android sets a limit at around 4-15
  - nRF52 s140 softdevice allows up to 20

- Connection management is often done in firmware
  - Softdevice for nRF, firmware on the radio chip in smartphones

- Limited by memory and complexity

How much throughput can a device have?

- 1 Mbps?

How much throughput can a device have?

- **1 Mbps?** Not even close. Packet overhead plus timing delays

- Step 1: decrease connection interval as much as possible
  - More connection events per second mean more data
  - Range: 7.5 ms to 4.0 s
  - Somewhat device-specific configuration
    - Android allows 7.5 ms
    - iOS allows 15 ms

How much throughput can a device have?

- Step 2, increase packet size

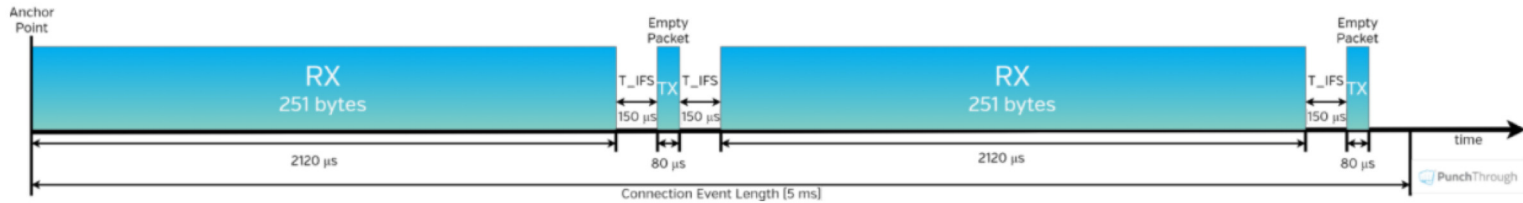Figure 5 – A single connection event (from the slave's perspective) in a DLE-enabled connection in which the master is transmitting as much data as possible within the effective Connection Event Length
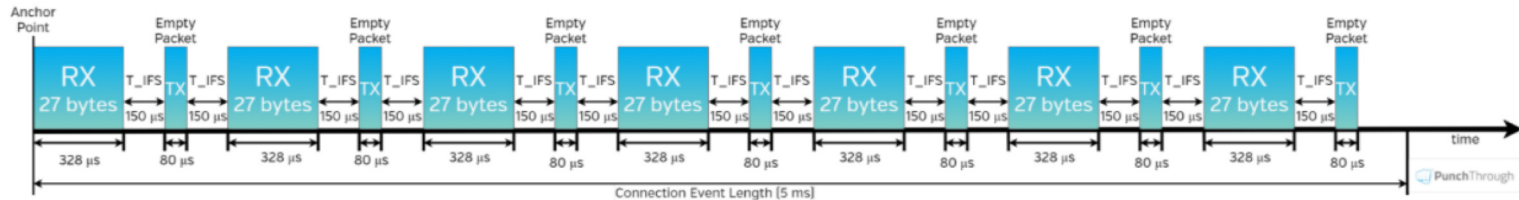
Figure 6 – A single connection event (from the slave's perspective) in a connection without DLE, in which the master is transmitting as much data as possible within the effective Connection Event Length

How much throughput can a device have?

- 488 bytes per connection event
  - Maximum sized packets, discounting headers and timing delays

- Connection event every 7.5 ms

- Result: 520 kbps (65 kB per second)
  - iOS result 260 kbps
  - Original BLE 4.1 result on Android: 128 kbps
  - Lower in practice due to lost packets

Outline

- Connection Link Layer

- Connections as Networks

- **GATT**

- BLE 5

Overview of Generic Attribute Profile (GATT)
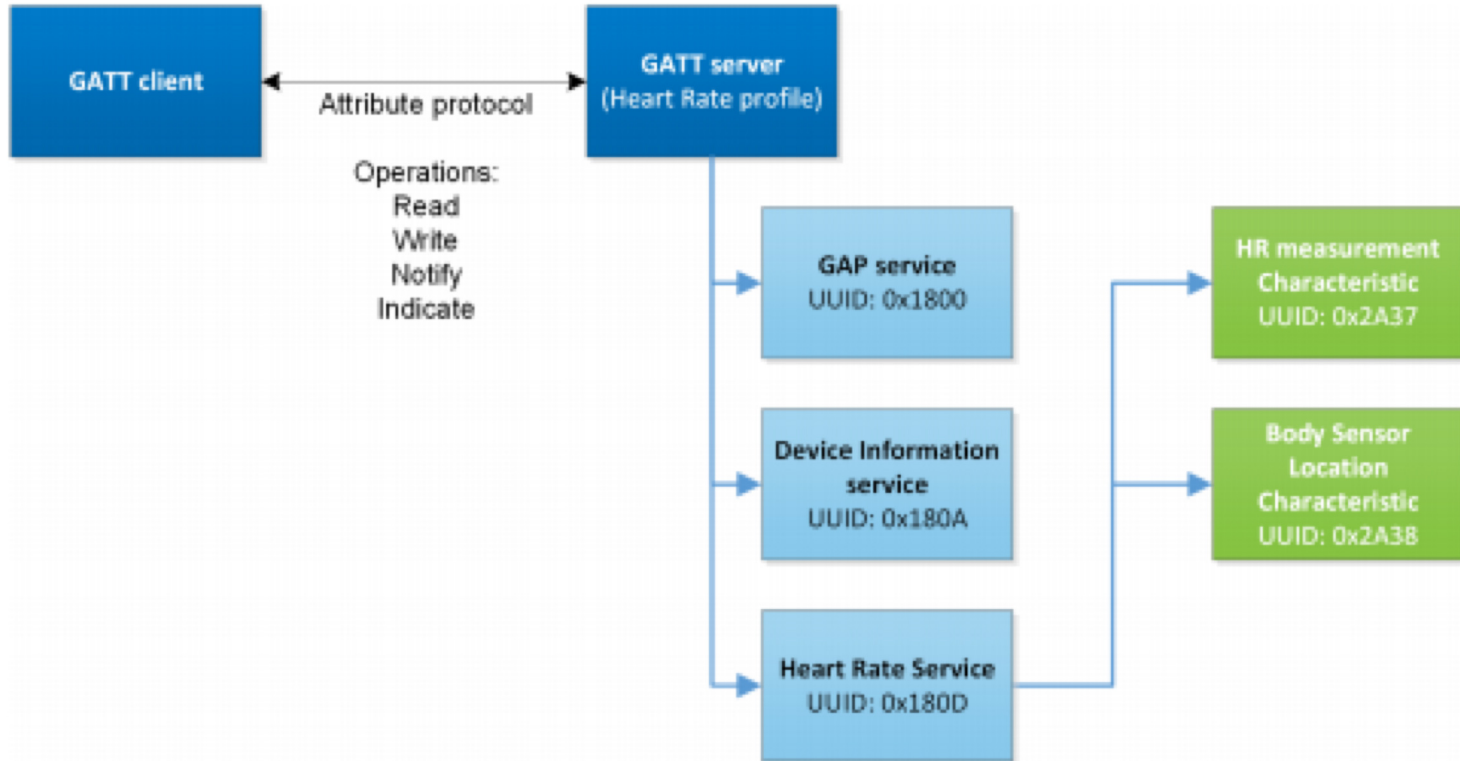
Attribute server keywords

- Characteristic
  - A field with properties and a value
  - Descriptor: metadata about the characteristic

- Service
  - Collection of characteristics

- Profile
  - Collection of services

Example: Time Profile

Current time characteristic

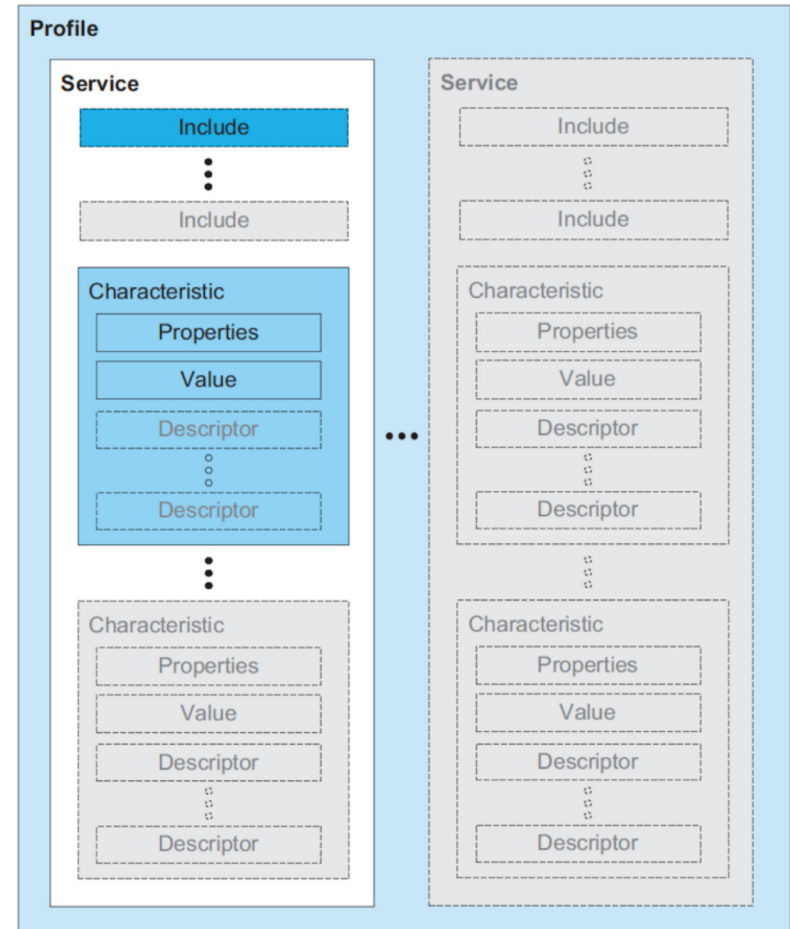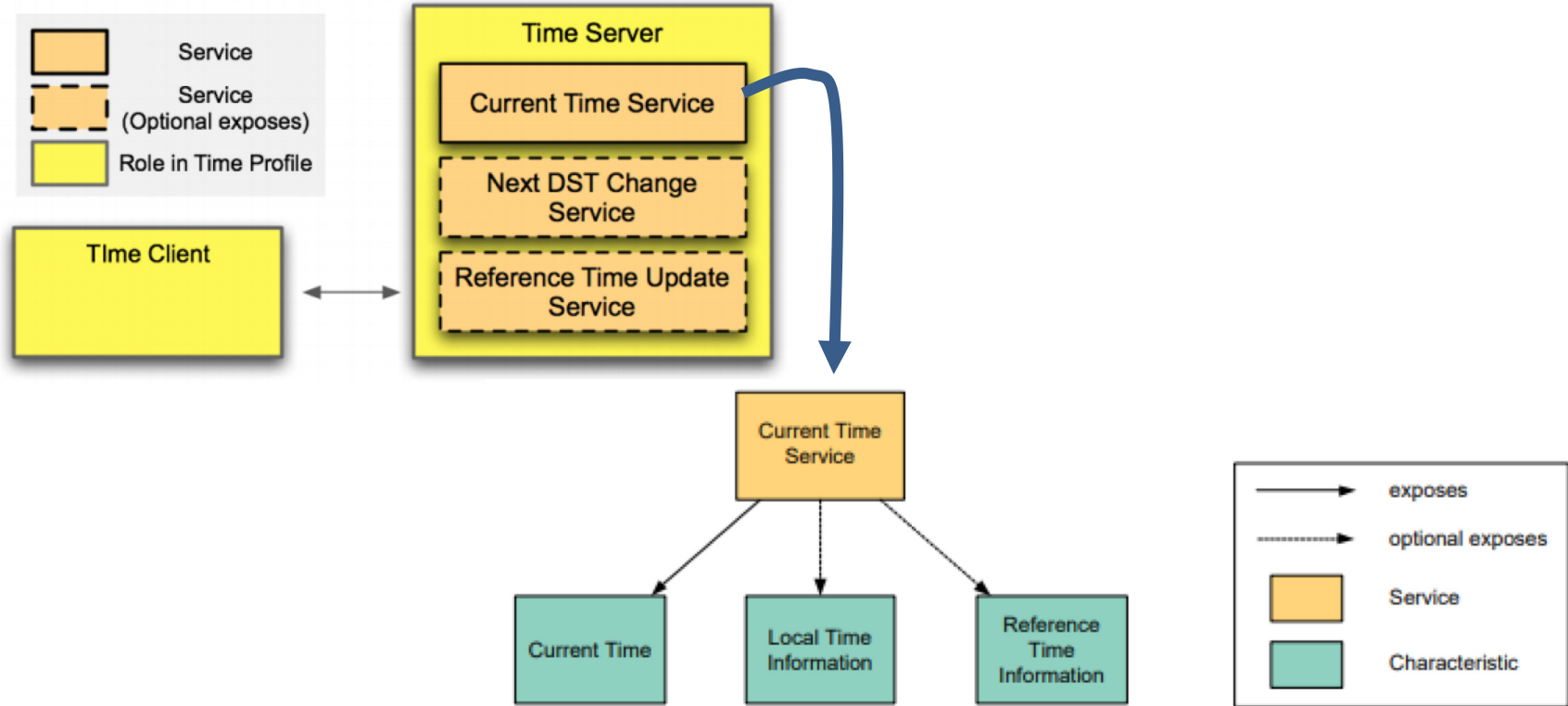| Field	Data Type	Size (in octets)	Description
Exact Time 256	struct	9	Refer to the Exact Time 256 characteristic in Section 3.64
Adjust Reason	uint8	1	See Section 3.58.2.1

Bit	Bit Name
0	Manual Time Update
1	External Reference Time Update
2	Change of Time Zone
3	Change of DST
4–7	Reserved for Future Use

Field	Data Type	Size (in octets)	Description
Day Date Time	struct	8	Refer to the Day Date Time characteristic in Section 3.54.
Fractions256	uint8	1	The number of 1/256 fractions of a second. Valid range 0–255.

Field	Data Type	Size (in octets)	Description
Date Time	struct	7	Refer to the Date Time characteristic in Section 3.53
Day of Week	struct	1	Refer to the Day of Week characteristic in Section 3.55

Field	Data Type	Size (in octets)	Description
Year	uint16	2	Year as defined by the Gregorian calendar. Valid range 1582 to 9999. A value of 0 means that the year is not known. All other values are reserved for future use (RFU).
Month	uint8	1	Month of the year as defined by the Gregorian calendar. Valid range 1 (January)

Documentation of GATT standards

- https://www.bluetooth.com/specifications/gatt/
  - Various profiles and services that have been standardized

- GATT Specification Supplement
  - Various characteristic definitions that have been standardized

- Both incredibly specific and woefully inexhaustive

UUIDs and handles

- Universally Unique Identifiers
  - 128-bit, mostly random with a few bits for versioning
  - Example: 00000000-0000-1000-8000-00805F9B34FB
    - This is the default BLE UUID for *known* services
    - You can generate your own UUID for custom services

- Handles
  - Too long to pass around all the time, so pick 16 bits that mean that UUID
    - Must be unique among services/characteristics on that device
  - Taken from UUID: 0000**xxxx**-0000-1000-8000-00805F9B34FB
    - 00001800-0000-1000-8000-00805F9B34FB   ← Generic Access Service [req'd]
    - 00002A05-0000-1000-8000-00805F9B34FB   ← Generic Attribute Service [opt]
  - Handle often sequentially incremented for each new characteristic within a service

Resource: Useful SO post on identifying handles in the wild

Discovery

- When a connection first occurs, each device can query the other for a list of services
  - And can further query for a list of characteristics in that service
  - Gets a list of handles/UUIDs

- Standardized UUIDs can be interpreted immediately
  - Custom services/characteristics need documentation from the manufacturer

Interacting with characteristics

- Depends on their permissions
  - Readable, Writable, Notify-able, etc.

- Notify
  - Automatically get a message sent whenever the characteristic value updates
  - Note: have to enable this on both sides, it's not the default behavior

- Long characteristics are automatically fragmented across multiple packets and/or connection events
  - L2CAP is in charge of this

Break + Example Services

Allocated UUID	Allocated for
0x1803	Link Loss
0x1804	Tx Power
0x1805	Current Time
0x1806	Reference Time Update
0x1807	Next DST Change
0x1808	Glucose
0x1809	Health Thermometer
0x180A	Device Information
0x180D	Heart Rate
0x180E	Phone Alert Status
0x180F	Battery
0x1810	Blood Pressure
0x1811	Alert Notification
0x1812	Human Interface Device
0x1813	Scan Parameters

Allocated UUID	Allocated for
0x1814	Running Speed and Cadence
0x1815	Automation IO
0x1816	Cycling Speed and Cadence
0x1818	Cycling Power
0x1819	Location and Navigation
0x181A	Environmental Sensing
0x181B	Body Composition
0x181C	User Data
0x181D	Weight Scale
0x181E	Bond Management
0x181F	Continuous Glucose Monitoring
0x1820	Internet Protocol Support
0x1821	Indoor Positioning
0x1822	Pulse Oximeter
0x1823	HTTP Proxy
0x1824	Transport Discovery
0x1825	Object Transfer
0x1826	Fitness Machine

Outline

- Connection Link Layer

- Connections as Networks

- GATT

- **BLE 5**

Changes in BLE 5

- Major changes
  - Multiple physical layers (optional implementation 😡)
  - Advertising extensions
  - Localization extensions (will discuss later with localization)

- Minor changes
  - Various quality of life improvements
  - Examples:
    - Advertise on channels in any order
   - Better data channel hopping algorithm

Revised physical layers

- 2 Mbps PHY
  - Transmit data faster
  - Transmit more data in the same time

- Coded PHY
  - Forward Error Correction in the data stream
    - 1 bit -> 2 symbols or 8 symbols
  - Makes bits more reliable -> longer distance
  - 500 kbps and 125 kbps modes

- Connections can switch to these PHYs after creation
- Advertisements can use these with extensions only

Coded PHY mixes physical and link layers

- Different PHY settings at different times
  - Make beginning headers extra-reliable
  - Data might be slightly less reliable as a trade for faster speed
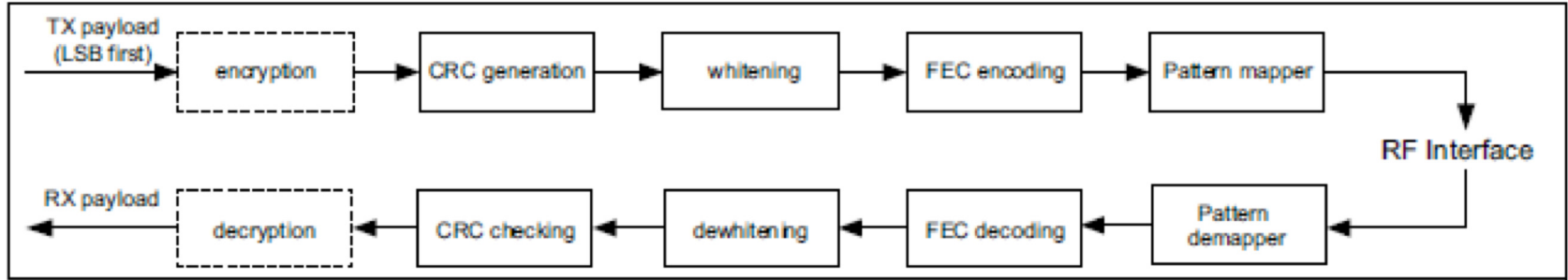
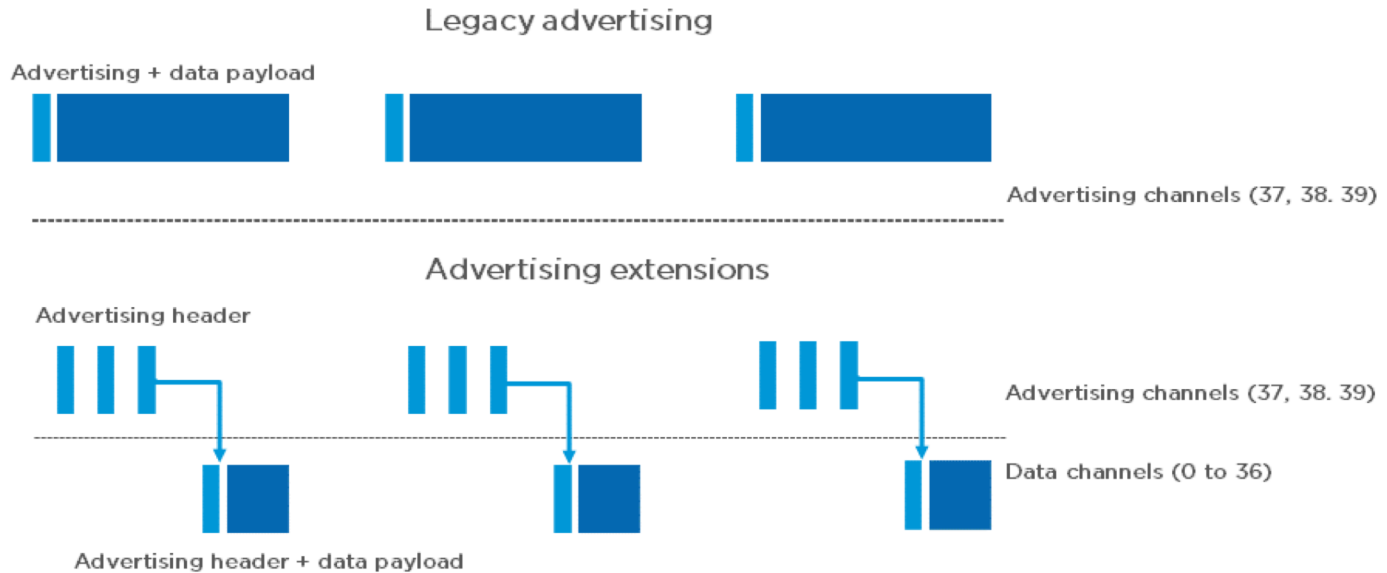Packet sent with coded PHY
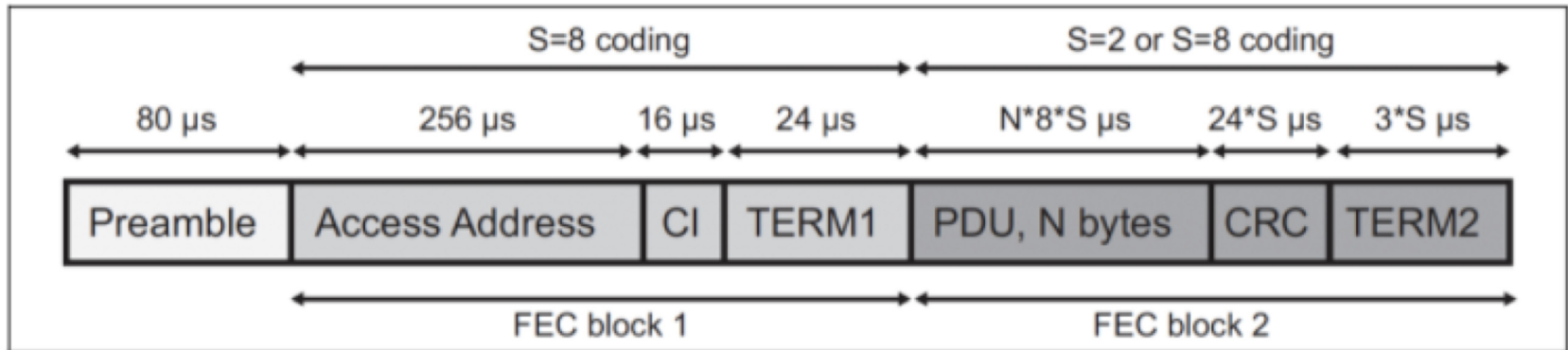
Revised processing path


Figure 3.2: Bit stream processing for the LE Coded PHYs

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Extended advertising

- Allow bigger payloads and/or different PHYs
  - Uses Data Channels to do so. Why?

-

Legacy advertising

Advertising + data payload

Advertising channels (37, 38. 39)

Advertising extensions

Advertising header

Advertising channels (37, 38. 39)

Data channels (0 to 36)

Advertising header + data payload

Extended advertising

- Allow bigger payloads and/or different PHYs
  - Uses Data Channels to do so. Why? Packet collisions!
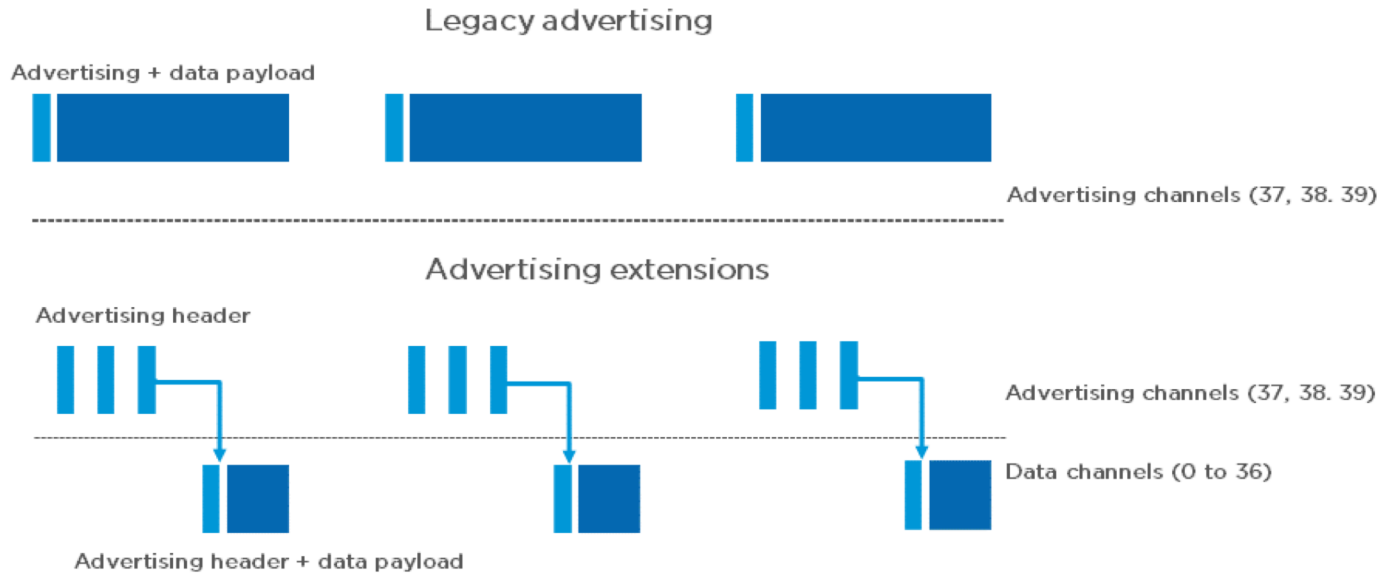
Procedure for scanning extended advertisements

1. Scan on 3 primary channels for advertising packets.
2. If ADV_EXT_IND is scanned, record the secondary channel information (which channel and when etc.)
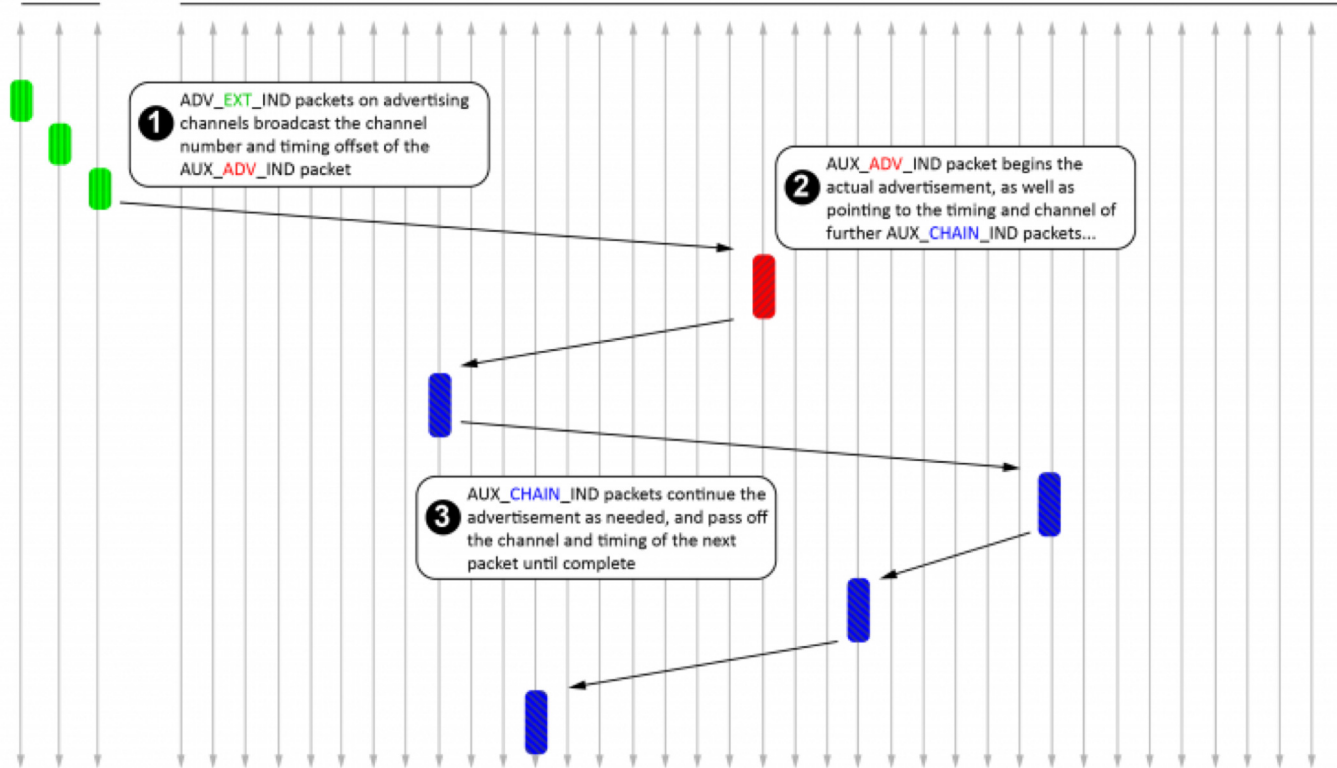3. Scan the specific secondary channel at the given time.

Extended advertisement train on data channels

Use case:
long advertisements
that need to be
fragmented

Periodic advertising on data channels
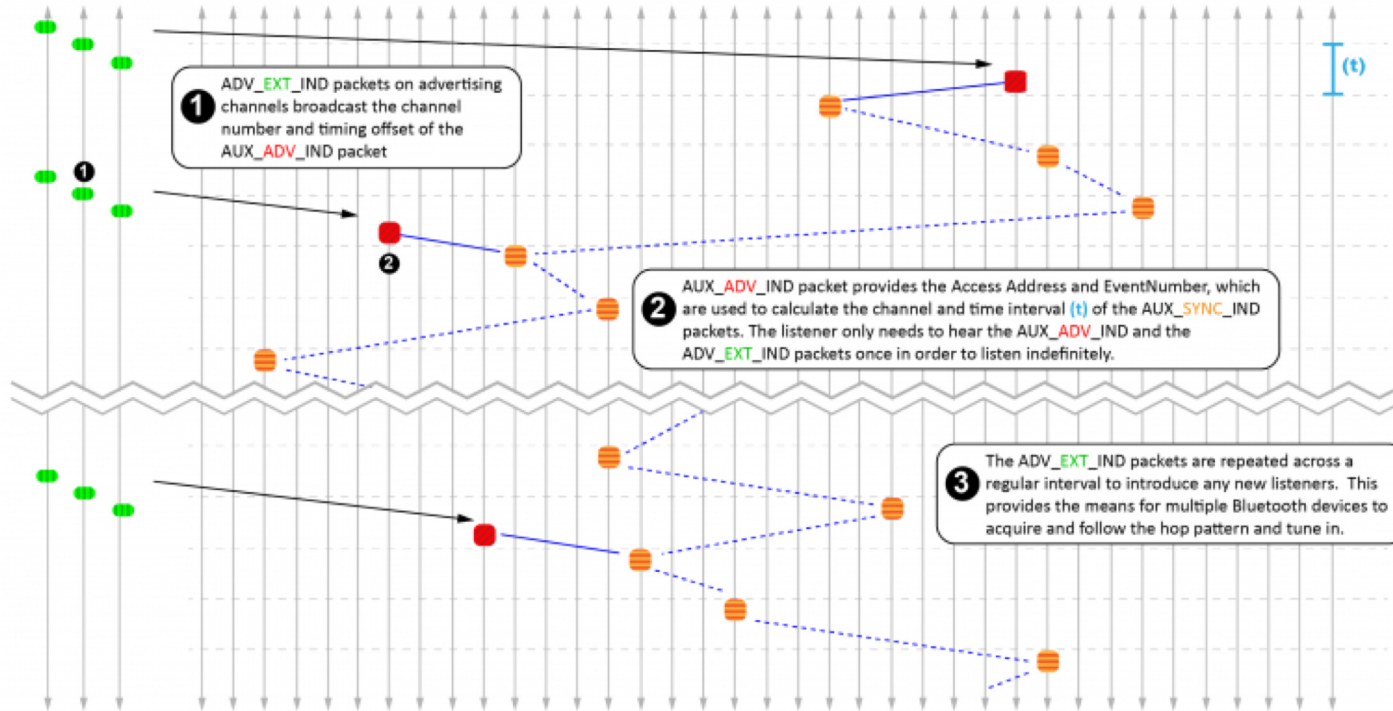
Use case: publicly-available localized audio sources

- TV in a bar
- Commentary in a museum

Broadcast analogy of a connection

Downside: how many of these can a gateway follow at a time?

Next Time: New PHY who dis?

BLE's cousin, IEEE 802.15.4

- Very similar modulation
  - In last ~5 years, many radio ICs are now `dual-mode`

- Very different history, use cases
  - Significantly more research, exploration
  - PAN vs WAN